

# Asymptotic bit frequency in Fibonacci words

JEAN-LUC BARIL  
LIB

Université de Bourgogne Franche-Comté  
B.P. 47 870, 21078 Dijon Cedex France  
email: barjl@u-bourgogne.fr

and  
SERGEY KIRGIZOV  
LIB

Université de Bourgogne Franche-Comté  
B.P. 47 870, 21078 Dijon Cedex France  
email: sergey.kirgizov@u-bourgogne.fr

and  
VINCENT VAJNOVSZKI  
LIB

Université de Bourgogne Franche-Comté  
B.P. 47 870, 21078 Dijon Cedex France  
email: vvajnov@u-bourgogne.fr

(Received: March 31, 2022, and in revised form May 15, 2022.)

**Abstract.** It is known that binary words containing no  $k$  consecutive 1s are enumerated by  $k$ -step Fibonacci numbers. In this note we discuss the expected value of a random bit in a random word of length  $n$  having this property.

**Mathematics Subject Classification(2020).** 05A16, 05A15, 11B39, 68R15.

**Keywords:** Fibonacci, word, pattern, frequency, asymptotics.

## 1 Introduction

For  $n \geq 0$  and  $k \geq 2$ , we denote by  $\mathcal{B}_n(1^k)$  the set of length  $n$  binary words avoiding  $k$  consecutive 1s. For example, we have

$$\begin{aligned}\mathcal{B}_4(11) &= \{0000, 0001, 0010, 0100, 0101, 1000, 1001, 1010\}, \text{ and} \\ \mathcal{B}_4(111) &= \{0000, 0001, 0010, 0011, 0100, 0101, 0110, 1000, 1001, 1010, 1011, 1100, 1101\}.\end{aligned}$$

It is well known, see Knuth [12, p. 286], that  $\mathcal{B}_n(1^k)$  is enumerated by the  $k$ -step Fibonacci numbers,

precisely  $|\mathcal{B}_n(1^k)| = f_{n+k,k}$ , where  $f_{n,k}$  is defined, following Miles [14] as

$$f_{n,k} = \begin{cases} 0 & \text{if } 0 \leq n \leq k-2, \\ 1 & \text{if } n = k-1, \\ \sum_{i=1}^k f_{n-i,k} & \text{otherwise.} \end{cases}$$

Denote by  $v_{n,k}$  the *frequency* (also called *popularity*) of 1s in  $\mathcal{B}_n(1^k)$ , i.e. the total number of 1s in all words of  $\mathcal{B}_n(1^k)$ . For instance,  $v_{4,2} = 10$  and  $v_{4,3} = 22$ . The ratio of frequency of 1s to the overall number of bits in words of  $\mathcal{B}_n(1^k)$  is

$$\alpha_{n,k} = \frac{v_{n,k}}{n \cdot |\mathcal{B}_n(1^k)|},$$

and it equals the expected value of a random bit in a random word from  $\mathcal{B}_n(1^k)$ . In [2], the authors left without proof the fact that, for any  $k \geq 2$ ,  $\lim_{n \rightarrow \infty} \alpha_{n,k}$  converges to a non-zero value as  $n$  grows. This note is devoted to proving this fact, which apart from its interest *en soi* has practical counterparts. Indeed, words in  $\mathcal{B}_n(1^k)$  play a critical role in some telecommunication frame synchronization protocols, see for example [1, 3, 5], or in particular Fibonacci-like interconnection networks [8].

Our discussion is based on the bivariate generating function

$$F_k(x, y) = \sum_{n=0}^{\infty} \sum_{m=0}^{n-\lfloor \frac{n}{k} \rfloor} a_{n,m} x^n y^m$$

whose coefficient  $a_{n,m}$  equals the number of words from  $\mathcal{B}_n(1^k)$  containing exactly  $m$  1s. For  $k = 2$  and  $k = 3$ , Table 1 presents some values of  $a_{n,m}$  for small  $n$  and  $m$ .

$m \setminus n$	1	2	3	4	5	6	7	8	9	$m \setminus n$	1	2	3	4	5	6	7	8	9
0	1	1	1	1	1	1	1	1	1	0	1	1	1	1	1	1	1	1	1
1	1	2	3	4	5	6	7	8	9	1	1	2	3	4	5	6	7	8	9
2		1	3	6	10	15	21	28		2		1	3	6	10	15	21	28	36
3			1	4	10	20	35			3			2	7	16	30	50	77	
4				1	5	15				4				1	6	19	45	90	
5					1					5					3	16	51		

Table 1: First few values of  $a_{n,m}$  for  $k = 2$  (left) and  $k = 3$ .

## 2 Main result

Proposition 2.1 gives the expression of the generating function  $F_k(x, y)$ . Even though this result is already obtained in [2], in order to make the paper self-contained we give an alternative proof of it. Then we calculate the generating functions for the frequency of 1s and for the overall number of bits in  $\mathcal{B}_n(1^k)$  by means of classic generating functions manipulations (Propositions 2.2). Applying

Theorem 4.1 from [16], after ensuring that its conditions are satisfied, we obtain the main result of this note, Theorem 2.6. The evolution of the random bit expectation for  $k = 2$  and  $k = 3$  is presented on Figure 1 for small values of  $n$ . And numerical estimations for the limit value ( $n \rightarrow \infty$ ) of the random bit expectation, for small values of  $k$  are given in Table 2.

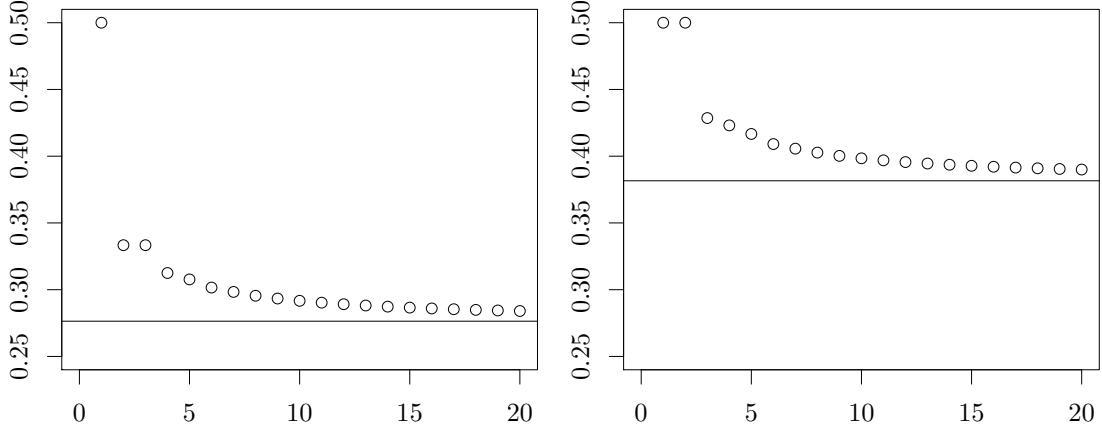


Figure 1: Expected value of a random bit in a random word from  $\mathcal{B}_n(1^2)$  (left) and  $\mathcal{B}_n(1^3)$  for small values of  $n$ .

PROPOSITION 2.1 ([2])

$$F_k(x, y) = \frac{y(1 - (xy)^k)}{y - xy^2 - xy + (xy)^{k+1}}.$$

Proof. The set  $\mathcal{B}(1^k) = \bigcup_{n=0}^{\infty} \mathcal{B}_n(1^k)$  respects the following recursive decomposition

$$\mathcal{B}(1^k) = \mathbb{1}_{k-1} \cup \left( \bigcup_{i=0}^{k-1} \left( 1^i 0 \cdot \mathcal{B}(1^k) \right) \right)$$

where  $\mathbb{1}_{k-1} = \bigcup_{i=0}^{k-1} \{1^i\}$  is the set of words in  $\mathcal{B}(1^k)$  containing no 0s, and  $\cdot$  denotes the concatenation. Note that the empty word also lies in  $\mathbb{1}_{k-1}$ . The claimed generating function is the solution of the following functional equation

$$F_k(x, y) = \sum_{i=0}^{k-1} x^i y^i + F_k(x, y) \sum_{i=0}^{k-1} x^{i+1} y^i.$$

□

In the proof of Theorem 2.6 we need the following easy to derive results.

PROPOSITION 2.2     •  $P_k(x) = \frac{\partial F_k(x, y)}{\partial y}|_{y=1}$  is the generating function where the coefficient of  $x^n$  is the frequency of 1s in  $\mathcal{B}_n(1^k)$ . We have

$$P_k(x) = \frac{x \cdot \sum_{i=0}^{k-2} (i+1)x^i}{(x^k + x^{k-1} + \dots + x^2 + x - 1)^2}.$$

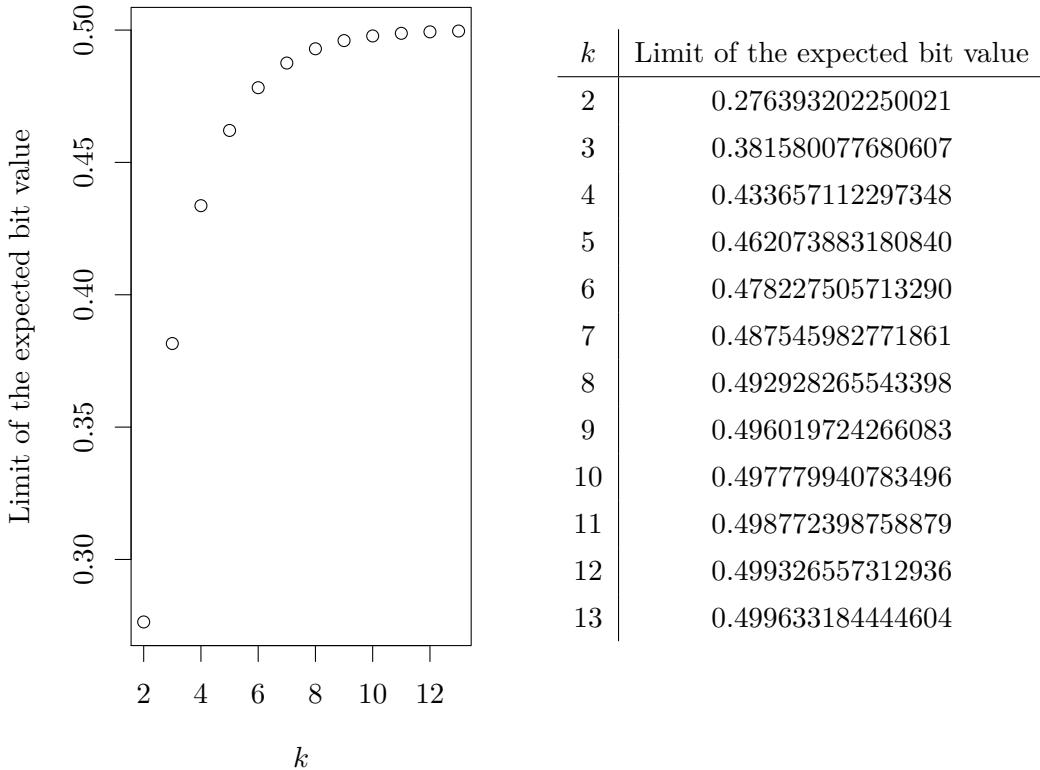


Table 2: Numerical estimations for the limit of the expected value of a random bit in a random word from  $\mathcal{B}_n(1^k)$ ,  $n \rightarrow \infty$ .

- $T_k(x) = x \frac{\partial F_k(x, 1)}{\partial x}$  is the generating function where the coefficient of  $x^n$  equals the total number of all bits in  $\mathcal{B}_n(1^k)$ . We have

$$T_k(x) = \frac{x \left( \sum_{i=0}^{k-2} (2i+2)x^i + \sum_{i=k-1}^{2k-2} (2k-i-1)x^i \right)}{(x^k + x^{k-1} + \cdots + x^2 + x - 1)^2}.$$

Every root  $r$  of a polynomial  $h(x)$  of degree  $n$  with a non-zero constant term corresponds to the root  $1/r$  of its negative reciprocal  $-x^n h(1/x)$ . The denominator of both  $P_k(x)$  and  $T_k(x)$  involves  $x^k + x^{k-1} + \cdots + x^2 + x - 1$  and its negative reciprocal is  $x^k - x^{k-1} - \cdots - x^2 - x - 1$  which is known in the literature as Fibonacci polynomial, see for instance [6, 7, 9, 10, 11, 13, 14, 15, 19] and references therein. In particular, Dubéau proved [7, Theorem 1] that its root of the largest modulus is  $\varphi_k = \lim_{n \rightarrow \infty} f_{n+1,k}/f_{n,k}$ , the generalized golden ratio, and  $\varphi_k$  approaches 2 when  $k \rightarrow \infty$  [7, Theorem 2]. Wolfram [19, Lemma 3.6] showed that any other root  $r$  of the Fibonacci polynomial satisfies

$3^{-1/k} < |r| < 1$ . See Figure 2 for an illustration of this fact. Moreover, Corollary 3.8 in [19] proves that Fibonacci polynomial is irreducible over  $\mathbb{Q}$ . In order to refer later to them we summarize these results in the next proposition.

PROPOSITION 2.3 *The polynomial  $g_k(x) = x^k + x^{k-1} + \cdots + x^2 + x - 1$  is irreducible over  $\mathbb{Q}$ , its root of the smallest modulus is unique and equal to  $1/\varphi_k$ .*

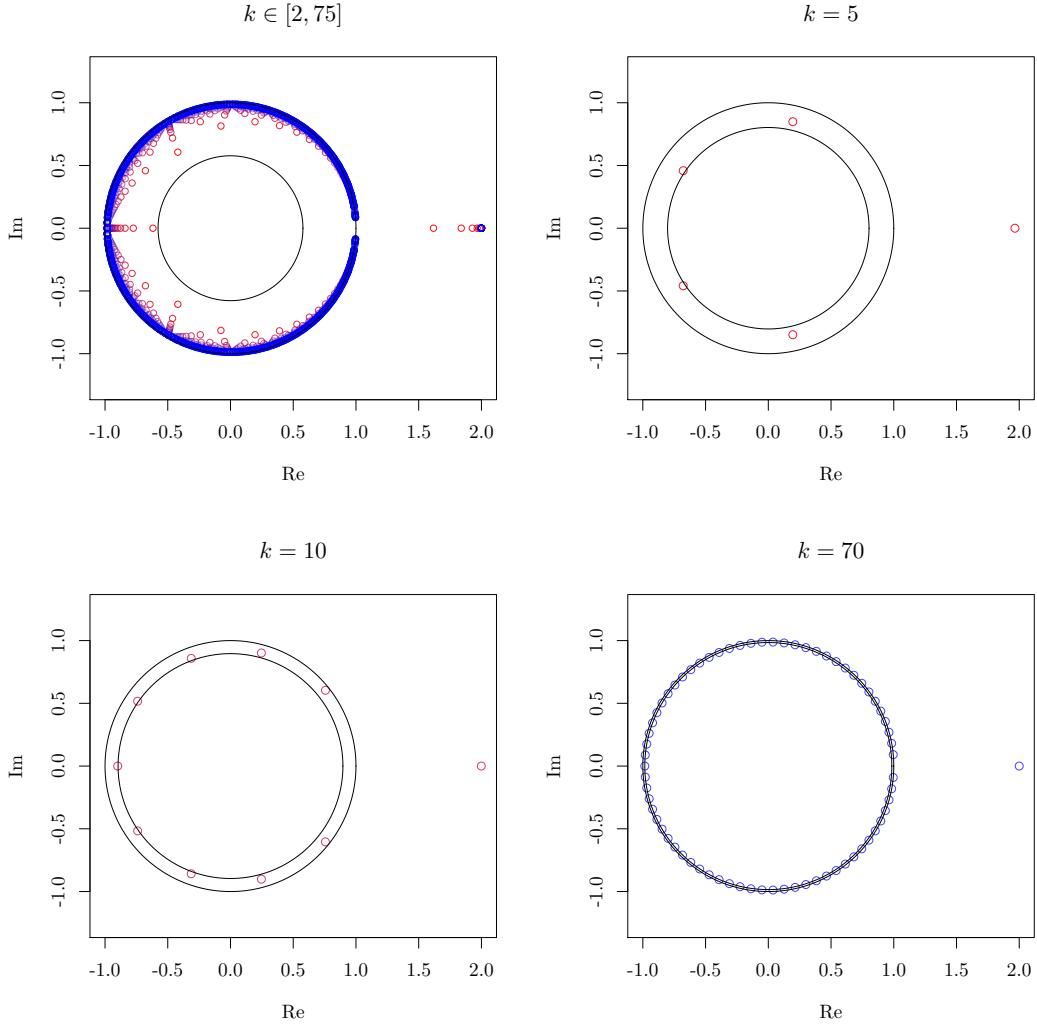


Figure 2: Roots of the polynomial  $x^k - x^{k-1} - \cdots - x^2 - x - 1$  (the negative reciprocal of  $g_k(x)$ ) for certain values of  $k$ .

The next lemma says that both fractions representing  $P_k(x)$  and  $T_k(x)$  are irreducible.

LEMMA 2.4 *The polynomials  $\sum_{i=0}^{k-2}(i+1)x^i$  and  $x^k + x^{k-1} + \cdots + x^2 + x - 1$  are relatively prime; and so are  $\sum_{i=0}^{k-2}(2i+2)x^i + \sum_{i=k-1}^{2k-2}(2k-i-1)x^i$  and  $x^k + x^{k-1} + \cdots + x^2 + x - 1$ .*

**Proof.** The polynomial  $x^k + x^{k-1} + \dots + x^2 + x - 1$  is irreducible due to Proposition 2.3. It does not divide  $\sum_{i=0}^{k-2} (i+1)x^i$  as it has a greater degree. And it also cannot divide  $\sum_{i=0}^{k-2} (2i+2)x^i + \sum_{i=k-1}^{2k-2} (2k-i-1)x^i$  as the latter does not have any positive real roots.  $\square$

From Propositions 2.2, 2.3, Dubeau's results [7], and Lemma 2.4 we have:

**LEMMA 2.5** *Both generating functions  $P_k(x)$  and of  $T_k(x)$  have the same and unique pole of the smallest modulus with multiplicity 2. The pole equals  $1/\varphi_k$ , where  $\varphi_k$  is the generalized golden ratio.*

For our main result of this note we need the Theorem 4.1 from [16]:

**Theorem 4.1 from [16].** *Assume that a rational generating function  $\frac{f(x)}{g(x)}$ , with  $f(x)$  and  $g(x)$  relatively prime and  $g(0) \neq 0$ , has a unique pole  $1/\beta$  of the smallest modulus. Then, if the multiplicity of  $1/\beta$  is  $\nu$ , we have*

$$[x^n] \frac{f(x)}{g(x)} \sim \nu \frac{(-\beta)^\nu f(1/\beta)}{g^{(\nu)}(1/\beta)} \beta^n n^{\nu-1}.$$

Both  $P_k(x)$  and  $T_k(x)$  are rational generating functions, and by Lemmas 2.4 and 2.5 they fulfill the conditions in the above theorem, so

$$\begin{aligned} [x^n] P_k(x) &\sim 2n\varphi_k^{n+2} \cdot \left. \frac{x \left( \sum_{i=0}^{k-2} (i+1)x^i \right)}{((x^k + x^{k-1} + \dots + x^2 + x - 1)^2)''} \right|_{x=1/\varphi_k} \\ [x^n] T_k(x) &\sim 2n\varphi_k^{n+2} \cdot \left. \frac{x \left( \sum_{i=0}^{k-2} (2i+2)x^i + \sum_{i=k-1}^{2k-2} (2k-i-1)x^i \right)}{((x^k + x^{k-1} + \dots + x^2 + x - 1)^2)''} \right|_{x=1/\varphi_k}. \end{aligned}$$

The expected value of a random bit in a random word from  $\mathcal{B}_n(1^k)$  is  $\frac{[x^n]P_k(x)}{[x^n]T_k(x)}$ . Taking the limit, we obtain:

**THEOREM 2.6** *The expected value of a random bit in a random word from  $\mathcal{B}_n(1^k)$  tends to*

$$\frac{kx^k - kx^{k-1} - x^k + 1}{kx^k - kx^{k-1} + x^{2k} - 3x^k + 2} \Big|_{x=1/\varphi_k} \quad \text{when } n \rightarrow \infty,$$

where  $\varphi_k = \lim_{n \rightarrow \infty} f_{n+1,k}/f_{n,k}$  is the generalized golden ratio, in particular  $\varphi_2$  is the golden ratio.

See Table 2 for some numerical estimations of the result obtained in the previous theorem. This result involves the generalized golden ratio. More than 20 years ago it was conjectured by Wolfram [19] that the Galois group of the polynomial  $x^k - x^{k-1} - \dots - x^2 - x - 1$  is the symmetric group  $S_k$ , and so there is no algebraic expression for  $\varphi_k$  (the root of the largest modulus of this polynomial) when  $k \geq 5$ . In case of even or prime  $k$  the conjecture was settled by Martin [13]. Cipu and Luca [6] showed that  $\varphi_k$  cannot be constructed by ruler and compass for  $k \geq 3$ . Nevertheless, good approximations are available, for instance Hare, Prodinger and Shallit [11] expressed  $\varphi_k$  and  $1/\varphi_k$  in terms of rapidly converging series.

The generalized golden ratio  $\varphi_k$  tends to 2 as  $k$  grows, and we deduce the following.

**COROLLARY 2.7** *The limit of the expected bit value of binary words avoiding  $k$  consecutive 1s, whose length tends to infinity, approaches  $1/2$  as  $k$  grows:*

$$\lim_{k \rightarrow \infty} \lim_{n \rightarrow \infty} \frac{v_{n,k}}{n \cdot |\mathcal{B}_n(1^k)|} = \frac{1}{2}.$$

Finally, note that other sets of restricted binary words are counted by the generalized Fibonacci numbers, for instance  $q$ -decreasing words [2] for  $q \geq 1$ . In this case every length maximal factor of the form  $0^a 1^b$  satisfies  $a = 0$  or  $q \cdot a > b$ . Theorem 2.6 and Corollary 2.7 apply to this case (with the same limit, see [2, Corollary 5]) by setting  $k = q + 1$ .

**Acknowledgement.** We would like to greatly thank Dietrich Burde, Ted Shifrin, Igor Rivin and Sil from Mathematics Stack Exchange<sup>1</sup> for insightful discussions and useful references about the irreducibility of Fibonacci polynomial which is directly related to Proposition 2.3. This work was supported in part by the project ANER ARTICO funded by Bourgogne-Franche-Comté region (France).

## References

- [1] D. BAJIC, *On construction of cross-bifix-free kernel sets*, 2nd MCM COST 2100, TD(07)237, Lisbon, Portugal, 2007.
- [2] J.-L BARIL, S. KIRGIZOV AND V. VAJNOVSZKI, *Gray codes for Fibonacci  $q$ -decreasing words*, arXiv:2010.09505.
- [3] A. BERNINI, S. BILOTTA, R. PINZANI AND V. VAJNOVSZKI, *A Gray code for cross-bifix-free sets*, Math. Structures Computer Sci., 27 (2017) 184–196.
- [4] A. BRAUER, *On algebraic equations with all but one root in the interior of the unit circle. To my teacher and former colleague Erhard Schmidt on his 75th birthday*, Math. Nachr., 4 (1950) 250–257.
- [5] Y. M. CHEE, H. M. KIAH, P. PURKAYASTHA AND C. WANG, *Cross-bifix-free codes within a constant factor of optimality*, IEEE Trans. Inform. Theory, 59 (2013) 4668—4674.
- [6] M. CIPU AND F. LUCA, *On the Galois group of the generalized Fibonacci polynomial*, An. Ştiinț. Univ. "Ovidius" Constanța Ser. Mat., 9 (2001) 27—38.
- [7] F. DUBEAU, *On  $r$ -generalized Fibonacci numbers*, Fibonacci Quart., 27 (1989) 221–229.
- [8] Ö. EĞECIOĞLU AND V. IRŠIČ, *Fibonacci-run graphs I: Basic properties*, Discrete Appl. Math., 295 (2021) 70–84.
- [9] I. FLORES, *Direct calculation of  $k$ -generalized Fibonacci numbers*, Fibonacci Quart., 5 (1967) 259–266.

<sup>1</sup>See the original discussions here <https://math.stackexchange.com/questions/4120185> and here <https://math.stackexchange.com/questions/4125568>.

- [10] G. W. GROSSMAN AND S. K. NARAYAN, *On the characteristic polynomial of the  $j$ -th order Fibonacci sequence*, Applications of Fibonacci Numbers, 8, Springer, Dordrecht, 1999, pp. 165–177.
- [11] K. HARE, H. PRODINGER AND J. SHALLIT, *Three series for the generalized golden mean*, Fibonacci Quart., 52 (2014) 307–314
- [12] D. KNUTH, *The Art of Computer Programming, Volume 3: Sorting and Searching*, 2nd ed., Addison-Wesley, 1998.
- [13] P. A. MARTIN, *The Galois group of  $x^n - x^{n-1} - \dots - x - 1$* , J. Pure Appl. Algebra, 190 (2004) 213–223.
- [14] E. MILES, *Generalized Fibonacci numbers and associated matrices*, Amer. Math. Monthly, 67 (1960) 745–752.
- [15] M. D. MILLER, *On generalized Fibonacci numbers*, Amer. Math. Monthly, 78 (1971) 1108–1109.
- [16] R. SEDGEWICK AND P. FLAJOLET, *An introduction to the analysis of algorithms*, 2nd ed., Addison-Wesley, 2013.
- [17] E. S. SELMER, *On the irreducibility of certain trinomials*, Math. Scand., 4 (1956) 287–302.
- [18] E. ROUCHÉ, *Mémoire sur la série de Lagrange*, Journal de l’École Polytechnique, 22 (1862) 193–224.
- [19] D. A. WOLFRAM, *Solving generalized Fibonacci recurrences*, Fibonacci Quart., 36 (1998) 129–145.